The Probabilistic Method - Probabilistic Techniques

Lecture 2: "The Method of Positive Probability (II)"

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The Basic Method (method of "positive probability")

- Construct (by using abstract random experiments) an appropriate probability sample space of combinatorial structures (thus, the sample points correspond to the combinatorial structures whose existence we try to prove).
- Prove that the probability of the desired property in this space is positive (i.e. non-zero).

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There is at least one point in the space with the desired property.

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There is at least one combinatorial structure with the desired property.

Examples in this lecture

- (i) Coloring Hypergraphs
- (ii) Tournaments with property S_k

Definition 1

A Hypergraph H = (V, E) consists of:

V : a finite set of vertices

E : a set of subsets of V (the "edges")

Definition 2

A Hypergraph H = (V, E) is called **n-uniform** iff all edges contain exactly n vertices.

Property B

Definition 3

A Hypergraph H = (V, E) has **property B** (it is two-colorable) iff \exists a two-coloring of V such that no edge is monochromatic.

Definition 4

m(n) is the minimum number of edges on a n-uniform hypergraph that **does not** have property B.

Theorem 1 (Erdös, 1963)

$$m(n) \ge 2^{n-1}$$

Proof of Theorem 1 (1/2)

- Construct a probability sample space by two-coloring the vertices of H at random, equiprobably for the two colors and independently for every vertex.
- Let *e* be any fixed edge.
- Define the event $M_e := \{e \text{ is monochromatic}\}.$
 - i.e. all vertices of edge *e* must have the same color.
- Compute the probability $\Pr[M_e]$.

$$\Pr[M_e] = 2 \cdot \underbrace{\frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} \cdots \frac{1}{2}}_{n \text{ times}} = 2^{1-n}$$

Proof of Theorem 1 (2/2)

- Define the event $M := \{ \exists \text{ at least one monochromatic edge} \}.$
- Hence, $M = \bigcup_e M_e$
- \blacksquare Using Boole's inequality we can compute $\Pr[M]$

$$\Pr[M] \leq \sum_e \Pr[M_e] = |e| 2^{1-n}$$

- If $|e| \cdot 2^{1-n} < 1$ (i.e., $|e| < 2^{n-1}$) then $\Pr[M] < 1 \Rightarrow \Pr[\overline{M}] > 0$.
- Hence, there is a two-coloring without a monochromatic edge when $m(n) < 2^{n-1} \Rightarrow$ property B.
- Hence, $m(n) \ge 2^{n-1}$ is necessary for avoiding property B.

Definition 5

A **tournament** T_n is a complete directed graph on n vertices i.e., for every pair (i, j), there is either an edge from i to j or from j to i, but not both.

Why do we call these graphs tournaments?

- Each vertex corresponds to a team playing at some tournament.
- The directed edge (i, j) means that team i wins team j.
- all teams play against each other.

Definition 6

A tournament T_n is said to have **property** S_k if for any set of k vertices in the tournament, there is some vertex that has a directed edge to each of those k vertices.

Theorem 2 (Erdös, 1963)

 $\forall k, \exists a \text{ tournament } T_n \text{ that has the property } S_k.$

Proof of Theorem 2 (1/2)

- Construct a probability sample space with points random tournaments by choosing the direction of each edge at random, equiprobably for the two directions and independently for every edge.
- Let S be any fixed set of k teams and define the event $M_S := \{ \nexists \text{ a team that wins all teams in } S \}.$
- For any team, the probability to win all teams in S is $(\frac{1}{2})^k$.
- Hence, the probability of not winning at least one of them is $1 (\frac{1}{2})^k$.
- The probability that this is happening for all *n* − *k* teams that don't belong in *S* is:

$$\Pr[M_S] = \left(1 - \left(\frac{1}{2}\right)^k\right)^{n-k}$$

Proof of Theorem 2 (2/2)

- Define the event M := {∃ a set S of k teams such that ∄ a team u : u ∉ S that wins all teams in S}.
- $\bullet \ M = \bigcup_S M_S$
- Using Boole's inequality we can compute Pr[M]

$$\Pr[M] \le \sum_{S, |S|=k} \Pr[M_S] = \binom{n}{k} \left(1 - \left(\frac{1}{2}\right)^k\right)^{n-k}$$

- If $\binom{n}{k} \left(1 \left(\frac{1}{2}\right)^k\right)^{n-k} < 1$ then $\Pr[M] < 1 \Rightarrow \Pr[\overline{M}] > 0$.
- Hence, there is a tournament with property S_k .