

# Παράλληλη Επεξεργασία

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«OpenMP - I»

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# Sequential Version

```
long num_steps = 100000;
double step;

int main()
{
    double x, pi, sum = 0.0;

    step = 1.0/(double) num_steps;

    for (int i=0; i <num_steps; i++){
        x = (i-0.5)*step;
        sum = sum + 4.0/(1.0+x*x);
    }

    pi = step * sum;

    return 0;
}
```

# POSIX Threads Version

```
#include <pthread.h>
#define NUM_THREADS 2
pthread_t thread[NUM_THREADS];
pthread_mutex_t Mutex;
long num_steps = 100000;
double step;
double global_sum = 0.0;

void *Pi (void *arg)
{
    int i, start;
    double x, sum = 0.0;

    start = *(int *) arg;
    step = 1.0/(double) num_steps;

    for(i=start; i<num_steps; i+=NUM_THREADS)
    {
        x = (i+0.5)*step;
        sum = sum + 4.0/(1.0+x*x);
    }
    pthread_mutex_lock (&Mutex);
    global_sum += sum;
    pthread_mutex_unlock(&Mutex);

    return 0;
}
```

```
int main ()
{
    double pi;

    int Arg[NUM_THREADS];

    for(int i=0; i<NUM_THREADS; i++)
        threadArg[i] = i;

    pthread_mutex_init(&Mutex, NULL);

    for (int i=0; i<NUM_THREADS; i++)
        pthread_create(&thread[i], NULL,
                      Pi, &Arg[i]);

    for (int i=0; i<NUM_THREADS; i++)
        pthread_join(thread[i], NULL);

    pi = global_sum * step;

    return 0;
}
```

# OpenMP version

```
#include <omp.h>
long num_steps = 100000;
double step;
#define NUM_THREADS 2
int main ()
{
    double x, pi, sum = 0.0;

    step = 1.0/(double) num_steps;
    omp_set_num_threads(NUM_THREADS);
    #pragma omp parallel for reduction(+:sum) private(x)
    for (int i=0; i<num_steps; i++){
        x = (i+0.5)*step;
        sum = sum + 4.0/(1.0+x*x);
    }

    pi = step * sum;

    return 0;
}
```

# Schedule and Goals

- OpenMP - part 1
  - study the basic features of OpenMP
  - able to understand and write OpenMP programs
- OpenMP - part 2
  - how OpenMP works
  - how to optimize OpenMP / parallel code
  - study and discuss more examples
- OpenMP – part 3
  - tasking model

*"We need to create learning situations where we ask students to practice program reading, to predict program execution, and to understand program idioms."*

Mark Guzdial, Communications of the ACM, Vol. 60 No. 6, Pages 10-11

# Example 1

- Identify and fix any issues in the following OpenMP codes

---

```
1 int A[N], B[N];
2 int auxdot = 0, dot = 0;
3
4 #pragma omp parallel
5 {
6     #pragma omp for
7     for (int i=0 ; i< N; i++ ){
8         auxdot += A[i]*B[i];
9     }
10
11     #pragma omp critical
12     dot += auxdot ;
13 }
```

---

---

```
1 #pragma omp parallel
2 {
3     if( omp_get_thread_num() % 2 ){
4 #pragma omp barrier
5
6         // ...
7     }
8 }
```

---

# Example 2

- Implement an equivalent version of the following code without using OpenMP worksharing

---

```
1 // double A[N];
2 // int i;
3
4 #pragma omp parallel for schedule(dynamic, 1)
5 for (i = 0; i < N; i++)
6 {
7     A[i] = work(i);
8 }
```

---

# Example 3

- Parallelize the following code using OpenMP

---

```
1 void compute_max_density()
2 {
3     // This routine finds the value of max density (max_rho) and
4     // its location (max_i, max_j) - there are no duplicate values
5     double max_rho;
6     int max_i, max_j;
7     max_rho = rho_[0];
8     max_i = 0;
9     max_j = 0;
10
11     for (int i = 0; i < N_; ++i)
12     for (int j = 0; j < N_; ++j)
13     {
14         if (rho_[i*N_ + j] > max_rho)
15         {
16             max_rho = rho_[i*N_ + j];
17             max_i = i;
18             max_j = j;
19         }
20     }
21 }
```



# Outline

- Introduction to OpenMP
- Parallel regions
- Worksharing constructs
  - loops, sections. single
- Combined parallel worksharing
- Data environment
- Synchronization
  - critical, atomic, barrier, master
- Library routines
- Environment variables
- Examples

# OpenMP

- OpenMP: An Application Program Interface (API) for writing multithreaded applications
  - simple, portable, widely supported standard
  - facilitates the development of multithreaded code in Fortran, C and C++
  - suitable for shared memory platforms
- Three primary components
  - compiler directives - instruct the compiler to generate multithreaded code
  - library calls
  - environment variables

# Evolution of OpenMP

Date	Version
Oct 1997	Fortran 1.0
Oct 1998	C/C++ 1.0
Nov 1999	Fortran 1.1
Nov 2000	Fortran 2.0
Mar 2002	C/C++ 2.0
May 2005	OpenMP 2.5
May 2008	OpenMP 3.0
Jul 2011	OpenMP 3.1
Jul 2013	OpenMP 4.0
Nov 2015	OpenMP 4.5

<http://computing.llnl.gov/tutorials/openMP/>

- OpenMP specifications at [www.openmp.org](http://www.openmp.org)
  - OpenMP 3.1 (2011): C/C++, Fortran and Examples
  - OpenMP 4.0 (2013): Examples in a separate PDF file

# Syntax Format

- Compiler directives
  - C/C++
    - `#pragma omp construct [clause [clause] ...]`
  - Fortran
    - `C$OMP construct [clause [clause] ... ]`
    - `!$OMP construct [clause [clause] ... ]`
    - `*$OMP construct [clause [clause] ... ]`
- Since we use directives, **no changes** need to be made to a program for a compiler that does not support OpenMP

# OpenMP Directive

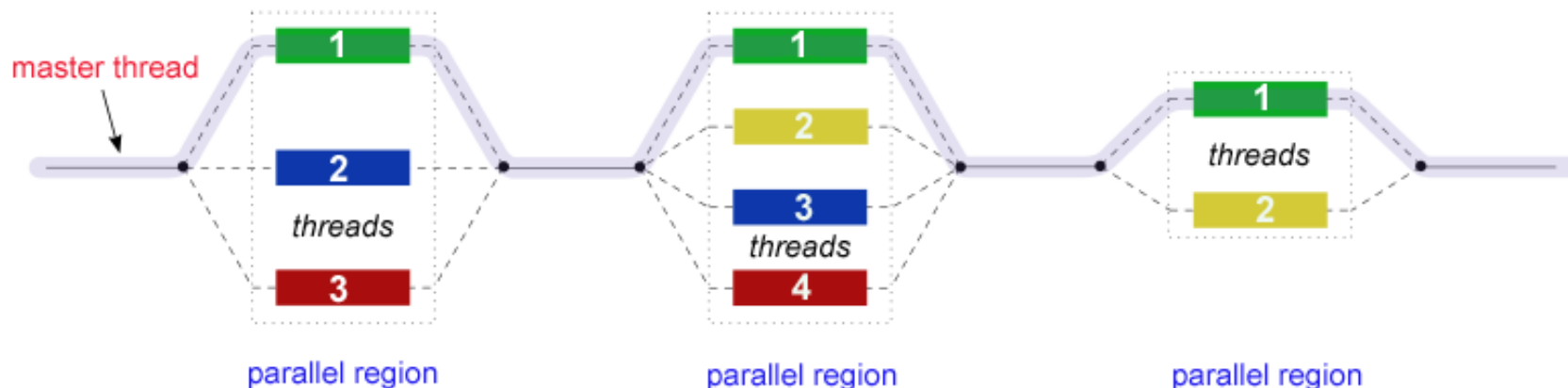
- Program executes serially until it encounters a parallel directive

```
#pragma omp parallel [clause list]  
/* structured block of code */
```

- Clause list is used to specify conditions
  - Conditional parallelism: **if (cond)**
  - Degree of concurrency: **num\_threads(int)**
  - Data handling: **private(vlist),**  
**firstprivate(vlist), shared(vlist)**

# Programming Model

- Fork-join type of parallelism:
  - The master thread spawns teams of threads according to the user / application requirements
  - Parallelism is added incrementally
    - the sequential code is transformed to parallel



<http://computing.llnl.gov/tutorials/openMP/>

# Typical Usage

- OpenMP is generally used for loop parallelization
  - Find the most time-consuming loops
  - Distribute the loop iterations to the threads

*Assign this loop to different threads*

```
void main()
{
    double Res[1000];
    for (int i=0;i<1000;i++) {
        do_huge_comp(Res[i]);
    }
}
```

Sequential code

```
void main()
{
    double Res[1000];
    #pragma omp parallel for
    for (int i=0;i<1000;i++) {
        do_huge_comp(Res[i]);
    }
}
```

Parallel code

But OpenMP is not just that!

# Using OpenMP

- Some compilers can automatically place directives with option
  - `-qsmp=auto` (IBM xlc)
  - some loops may speed up, some may slow down
- Compiler option required when you use directives
  - `-fopenmp` (GNU compilers)
  - `-openmp` (Intel compilers)
  - `-qsmp=omp` (IBM)
- Scoping variables can be sometimes the hard part!
  - shared variables, thread private variables



# Hello World!

```
#include <omp.h>
#include <stdio.h>
```

OpenMP include file

```
int main() {
    #pragma omp parallel
```

Parallel region with default  
number of threads

```
{
    int me = omp_get_thread_num();
    int nthr = omp_get_num_threads();
    printf("Hello world from thread %d of %d\n", me, nthr);
}
```

Library calls

```
return 0;
```

End of parallel region

```
}
```

- Compilation with the GNU GCC and Intel compilers

```
$ gcc -fopenmp -o hello hello.c
```

```
$ icc -openmp -o hello hello.c
```

MacOS: brew install gcc

# Usage

- Execution

```
$ export OMP_NUM_THREADS=4
```

```
$ ./hello
```

```
Hello world from thread 0 of 4
```

```
Hello world from thread 2 of 4
```

```
Hello world from thread 1 of 4
```

```
Hello world from thread 3 of 4
```

```
$ export OMP_NUM_THREADS=1
```

```
$ ./hello
```

```
Hello world from thread 0 of 1
```

Environment variable

# Thread Interaction

- OpenMP is a shared-memory programming model
  - Threads communicate through shared variables
- Data sharing can lead to **race conditions**
  - the output of some code can change due to thread scheduling, e.g. their order of execution
- Synchronization at the right places can eliminate race conditions
  - However, **synchronization is expensive**
  - the way data is stored might need to change to minimize the need for synchronization

# OpenMP Directives

- 5 categories
  - Parallel Regions
  - Worksharing
  - Data Environment
  - Synchronization
  - Runtime functions & environment variables
- Basically the same between C/C++ and Fortran

# Parallel Regions

- Create threads with **omp parallel**
- The following code will create a parallel region of 4 threads:

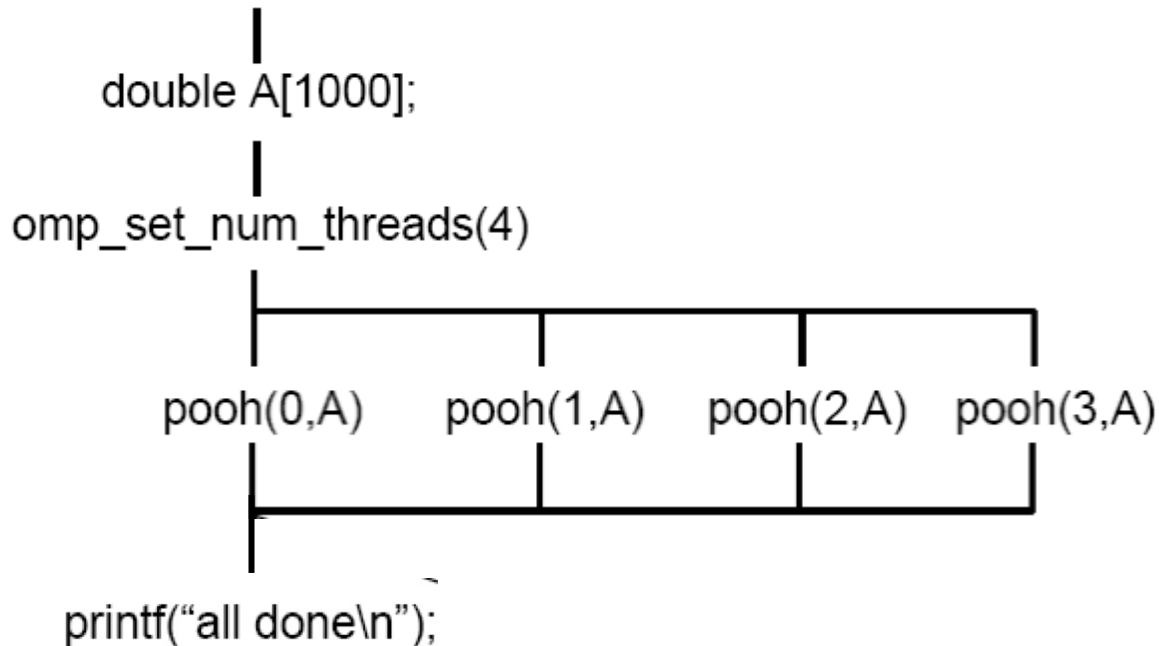
```
double A[1000];
omp_set_num_threads(4);
#pragma omp parallel
{
    int ID = omp_get_thread_num();
    pooh(ID,A);
}
```

- Threads share A (default behavior)
- Master thread creates the threads
- Threads all start at same time then synchronize at a barrier at the end to continue with code
- Each threads calls pooh for its own ID (0 to 3)

# Parallel Regions

- Each threads runs the same code
- All threads share A
- Execution continues when all threads have finished their work (barrier)

```
double A[1000];  
omp_set_num_threads(4);  
#pragma omp parallel  
{  
    int ID = omp_get_thread_num();  
    pooh(ID,A);  
}  
printf("all done\n");
```



# Parallel Regions - Syntax

```
#pragma omp parallel [clause ...] newline  
  
    structured_block
```

## Clauses

```
if (scalar_expression)  
num_threads (integer-expression)  
private (list)  
shared (list)  
firstprivate (list)  
default (shared | none)  
reduction (operator: list)  
copyin (list)
```

# Structured Blocks

- Most OpenMP directives are applied to structured blocks of code
  - Structured block: piece of code with a single entry point at the beginning and a single exit point at the end.

```
#pragma omp parallel
{
    int id =omp_get_thread_num();
    res[id] = work(id);
}
printf("after parallel\n");
```

Structured block

```
#pragma omp parallel
{
    int id = omp_get_thread_num();
    res[id] = work(id);
    if (res[id] == 0) goto out;
}
out: printf("after parallel\n");
```

Unstructured block



# Clauses for omp parallel

<code>if (scalar_expression)</code>	Only parallelize if the expression is true. Can be used to stop parallelization if the work is too little
<code>num_threads (integer-expression)</code>	Set the number of threads
<code>private (list)</code>	The specified variables are thread-private
<code>shared (list)</code>	The specified variables are shared among all threads
<code>firstprivate (list)</code>	The specified variables are thread-private and initialized from the master thread
<code>reduction (operator: list)</code>	Perform a reduction on the thread-local variables and assign it to the master thread
<code>default (shared   none)</code>	Unspecified variables are shared or not

```
#pragma omp parallel private(i) shared(n) if(n > 10)  
{  
    //...  
}
```

# Actual Number of Threads

- The number of threads in a parallel region is determined by the following factors, in order of precedence:
  1. Evaluation of the **if** clause
  2. Setting of the **num\_threads** clause
  3. Use of the **omp\_set\_num\_threads()** library function
  4. Setting of the **OMP\_NUM\_THREADS** environment variable
  5. Implementation default - usually the number of CPUs on a node, though it could be dynamic.
- Reminder: threads are numbered from 0 (master thread) to N-1

# Static and Dynamic modes

- Dynamic mode (default):
  - The number of threads can differ between parallel regions of the same program
  - The specified number of threads actually defines the maximum number - the actual number of threads can be smaller
- Static mode:
  - The number of threads is fixed and exactly equal to the number specified by the programmer
- OpenMP supports nested parallel regions but...
  - The compiler is allowed to serialize all the inner levels
  - This means that it uses a single OpenMP thread for those parallel regions

# Worksharing Constructs

- the for construct splits up loop iterations

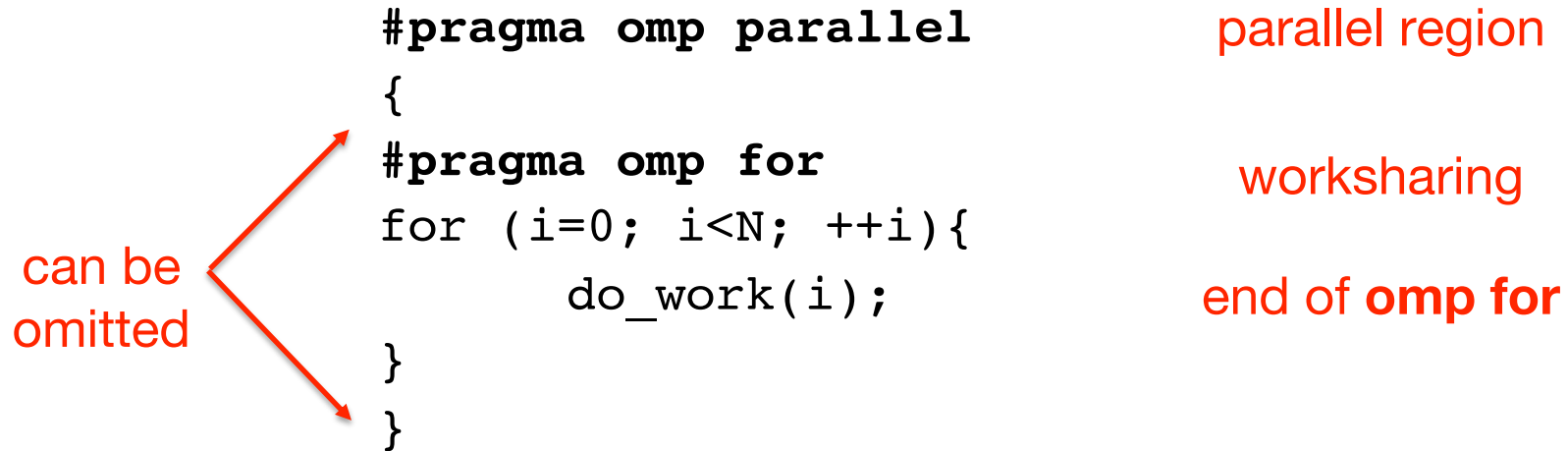
can be omitted

```
#pragma omp parallel
{
  #pragma omp for
  for (i=0; i<N; ++i){
    do_work(i);
  }
}
```

parallel region

worksharing

end of **omp for**

A diagram illustrating the components of an OpenMP worksharing construct. On the left, the text "can be omitted" is written in red. Two red arrows originate from this text: one points to the opening curly brace of the `#pragma omp parallel` block, and the other points to the closing curly brace of the `#pragma omp for` block. On the right, three red annotations are present: "parallel region" is aligned with the `#pragma omp parallel` line, "worksharing" is aligned with the `#pragma omp for` line, and "end of **omp for**" is aligned with the closing curly brace of the `for` loop.

- By default, there is a barrier at the end of the **omp for**.
- Use the **nowait** clause to turn off the barrier.

# Rule

- In order to be made parallel, a loop must have canonical “shape”

```
for (index=start; index < end; index++)
for (index=start; index <= end; ++index)
for (index=start; index >= end; index--)
for (index=start; index > end; --index)
for (index=start; index <= end; index += inc)
for (index=start; index >= end; index -= inc)
for (index=start; index < end; index = index + inc)
for (index=start; index < end; index = inc + index)
for (index=start; index < end; index = index - inc)
```

# Sections construct

- The sections construct gives a different structured block to each thread

```
#pragma omp parallel
#pragma omp sections
{
#pragma omp section
    x_calculation();
#pragma omp section
    y_calculation();
#pragma omp section
    z_calculation();
}
```

parallel region  
worksharing

each section gets assigned  
to a different thread

end of omp sections

- By default there is a barrier at the end. The **nowait** clause turns it off

# Single construct

- The structured block is executed only by one of the threads
- An implicit barrier exists at the end of **single**
- Can be considered as a synchronization construct

```
#pragma omp parallel
{
    do_many_things();
    #pragma omp single
    {
        exchange_boundaries();
    }
    do_many_other_things();
}
```

implicit barrier here

and here, end of parallel region

# Combined Directives

- Parallel regions can be combined with the `for` and `sections` worksharing constructs
- **`omp parallel + omp for` → `omp parallel for`**

```
#pragma omp parallel for  
for (i=0; i<N; i++){  
    do_work(i);  
}
```



# Combined Directives

- **omp parallel + omp sections →  
omp parallel sections**

```
#pragma omp parallel sections
{
  #pragma omp section
    x_calculation();
  #pragma omp section
    y_calculation();
  #pragma omp section
    z_calculation();
}
```

# Directive Scoping

- OpenMP directives can be extended in multiple files
- Orphan directives: appear outside a parallel region

```
//foo.c
```

```
#pragma omp parallel
```

```
{  
    whoami();  
}
```

```
//bar.c
```

```
void whoami()
```

```
{  
    int iam = omp_get_thread_num();  
    #pragma omp critical synchronization  
    {  
        printf("Hello from %d"\n, iam);  
    }  
    return;  
}
```

- foo.c: Static (lexical) extent of parallel region
- bar.c: Dynamic extent of parallel region

# Data Scoping

- OpenMP is a shared memory programming model
  - most variables are shared by default
- Global variables are shared
- But not everything is shared
  - loop index variables
  - stack variables in called functions from parallel region

# Storage Attributes

- The programmer can change the storage attributes of variables with the following clauses
  - **shared**
  - **private**
  - **firstprivate**
  - **threadprivate**
- The value of a private variable used in a parallel loop can be exported as global value with the clause:
  - **lastprivate**
- The default behavior can be changed using:
  - **default(private | shared | none)**
- The data clauses are applied to the parallel region and worksharing constructs - however, shared is only valid for parallel regions
- Data scoping clauses are valid only in the lexical extent of the OpenMP directive

# Data Environment

- Example of **private** and **firstprivate**

```
int A, B, C;  
A = B = C = 1;  
#pragma omp parallel private(B) firstprivate(C)  
{  
    // ...  
}
```

- Within the parallel region :
  - “A” is shared between threads and equal to 1
  - Both “B” and “C” are private for each thread
    - B has undefined initial value
    - C has initial value equal to 1
- After the parallel region:
  - Both B and C have the same value as before the parallel region

# private

- **private**(var) creates a private copy of var in each thread
  - The value of the copy is not initialized
  - The private copy is not related to the original variable with respect to the memory location

```
int is = 0;
#pragma omp parallel for private(is)
for (int j=1; j<=1000; j++)
    is = is + j;

printf("%d\n", is);
```

- IS has not been initialized inside the loop

# firstprivate

- **firstprivate**: special case of **private**
  - The private copy of each thread is initialized with the value of the original variable, which belongs to the master thread

```
int is = 0;
#pragma omp parallel for firstprivate(is)
for (int j=1; j<=1000; j++)
    is = is + j;

printf("%d\n", is);
```

- Each thread has a private copy of IS with initial value 0

# lastprivate

- Copies the value of the private variable, as assigned by the last loop iteration, to the original (global) variable

```
int is = 0;
#pragma omp parallel for firstprivate(is) \
lastprivate(is)
for (int j=1; j<=1000; j++)
    is = is + j;

printf("%d\n", is);
```

continue to  
the next line

- Each thread has a private copy of IS with initial value 0
- IS has the value it was assigned by the last loop iteration (i.e. for j=1000)



# Synchronization

- OpenMP supports several synchronization constructs:
  - `critical` section
  - `atomic`
  - `barrier`
  - `master` (in fact, not a synchronization construction )
  - *ordered* not studied
  - *flush* not studied

# Synchronization – critical

- No two threads will simultaneously be in the critical section
- Critical sections can be named
  - `omp critical (name)`

```
#pragma omp parallel for private(b) shared(res)
for (i=0; i<niters; i++) {
    b = doit(i);
    #pragma omp critical
    {
        update(b, &res);
    }
}
```

lock mutex

unlock mutex

res: initialized before the parallel region

# Synchronization – atomic

- Special case of critical section that can be used only for simple instructions.
- Can be applied only when a single memory location (variable) is updated

```
#pragma omp parallel private(b)
{
    int i = omp_get_thread_num();
    b = doit(i);
    #pragma omp atomic
        res = res + b;
}
```

use of some  
hardware-supported  
atomic operation

res: initialized before the parallel region

# Synchronization – barrier

- Barrier: all threads wait until each thread has reached the barrier

```
#pragma omp parallel shared (A, B) private(id)
```

```
{
```

```
    id=omp_get_thread_num();
```

```
    A[id] = big_calc1(id);
```

```
    #pragma omp barrier
```

initialization of A

necessary synchronization

```
    #pragma omp for
```

```
    for(int i=0; i<N; i++){
```

```
        B[i]=big_calc2(i,A);
```

```
    }
```

```
}
```

these computations  
depend on A

# Synchronization – master

- The structured block is executed only by the **master thread**  
- the other threads of the team ignore it
- There is no barrier at the end of master

```
#pragma omp parallel
{
    do_many_things();
    #pragma omp master
    {
        exchange_boundaries();
    }
    #pragma barrier
    do_many_other_things();
}
```

nothing more than  
if (omp\_get\_thread\_num()==0)

# Synchronization - Implicit Barriers

- A barrier is implicitly called at the end of the following constructs:
  - **parallel**
  - **for** (except when **nowait** is used)
  - **sections** (except when **nowait** is used)
  - **single** (except when **nowait** is used)
- **for**, **sections** and **single** accept the **nowait** clause

```
int nthreads;  
  
#pragma omp parallel  
#pragma omp single nowait  
nthreads = omp_get_num_threads();
```

# Reductions

- The reduction clause modifies the way variables are “shared”:
  - **reduction (op : list)**
- Variables included in **list** must be shared in the parallel region where the reduction clause exists
- Allowed reduction operations: `+, -, *, &, ^, |, &&, | |, min, max`
- Within a parallel region or a worksharing construct:
  - A local copy for each variable in the list is created and initialized accordingly to the reduction operation
    - 0 for “+”
  - The values of the local copies are combined (reduced) to a single value that is stored to the original variable after the end of the construct

# Reduction - Example

```
#include <omp.h>
#define NUM_THREADS 2

double func(int i);

void main ()
{
    int i;
    double ZZ, res=0.0;
    omp_set_num_threads(NUM_THREADS);
#pragma omp parallel for reduction(+:res) private(ZZ)
    for (i=0; i< 1000; i++){
        ZZ = func(i);
        res = res + ZZ;
    }
}
```



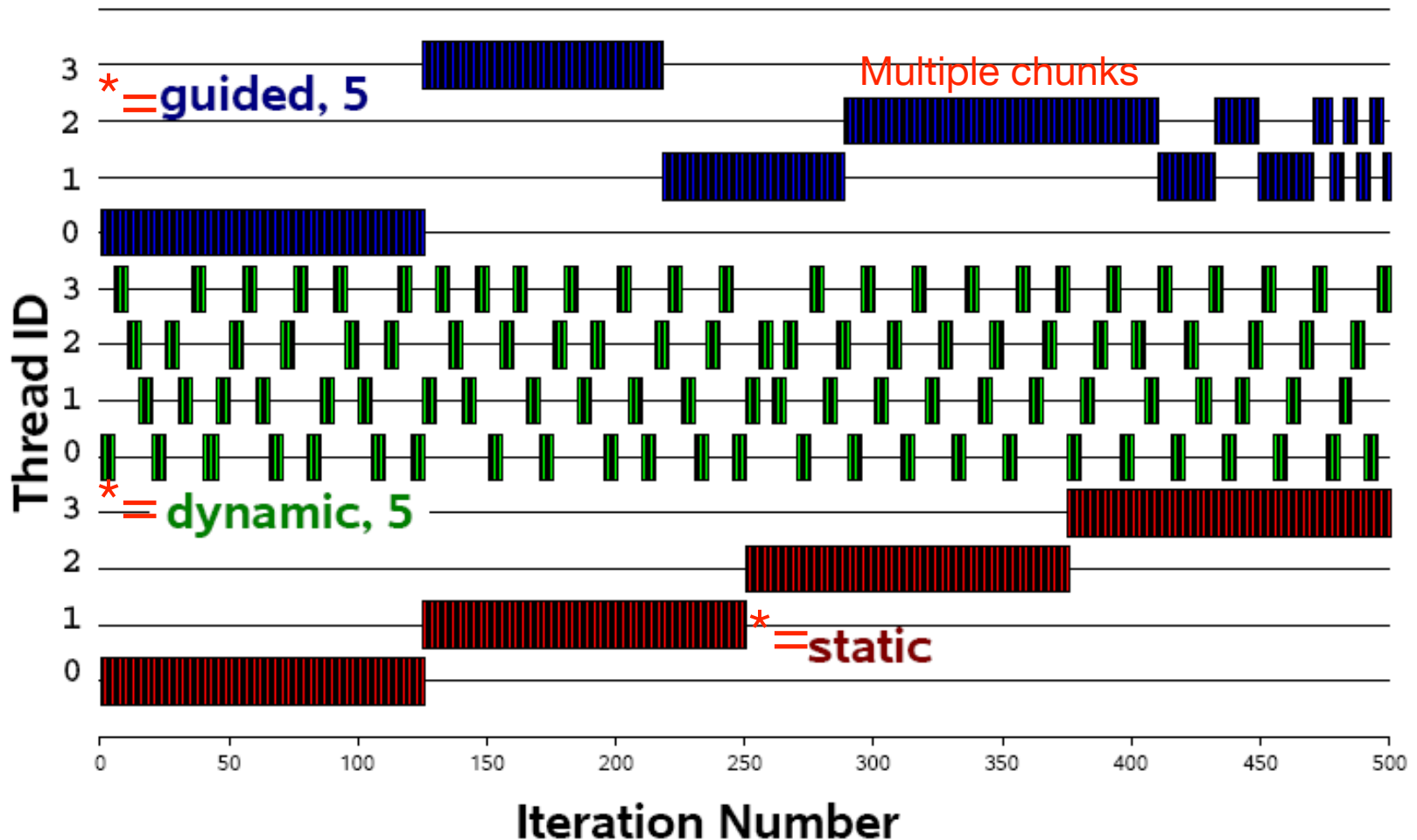
# Loop Scheduling

- Usage: `#pragma omp parallel for <schedule clause>`
  - `schedule ( static | dynamic | guided [, chunk] )`
  - `schedule (runtime)`
- **static** [, **chunk**]
  - Loop iterations are divided into segments of size **chunk** and distributed cyclically to the threads of the parallel region
  - If **chunk** is not specified, it is equal to  $N/P$  and each thread executes a single chunk of iterations
- **dynamic** [, **chunk**]
  - Loop iterations are divided into segments of size **chunk**
  - An idle thread gets dynamically the next available chunk of iterations
- **guided** [, **chunk**]
  - Similar to dynamic but the chunk size decreases exponentially.
  - **chunk** specifies the minimum segment size
- **runtime**
  - decide at runtime depending on the `OMP_SCHEDULE` environment variable
- **auto**
  - decided by the compiler and/or the underlying OpenMP runtime library

# Example

```
#pragma omp parallel for num_threads(4) schedule(*)  
for (int i = 0; i < 500; i++) do_work(i);
```

*500 iterations on 4 threads*



# Library Calls

- OpenMP locks
  - `omp_init_lock()`, `omp_set_lock()`, `omp_unset_lock()`,  
`omp_test_lock()`
- Functions that control the runtime environment:
  - Number of threads
    - `omp_set_num_threads()`, `omp_get_num_threads()`,  
`omp_get_thread_num()`, `omp_get_max_threads()`
  - Dynamic mode and nested parallelism
    - `omp_set_dynamic()`, `omp_set_nested()`,
    - `omp_get_dynamic()`, `omp_get_nested()`
  - Check if code is in a parallel region
    - `omp_in_parallel()`
  - Number of processors / cores
    - `omp_get_num_procs()`
- Wall-clock time measurement (in seconds)
  - `omp_get_wtime()`

ex01: `get_wtime()`

# OpenMP Locks

```
omp_lock_t lck;  
omp_init_lock(&lck);  
#pragma omp parallel
```

lock variable  
initialization

```
{  
    int id = omp_get_thread_num();  
    int tmp = do_lots_of_work(id);  
  
    omp_set_lock(&lck);  
    printf(" %d %d\n", id, tmp);  
    omp_unset_lock(&lck);  
}
```

```
omp_destroy_lock(&lck);
```

destruction

# Libraries Calls

- Dynamic mode is disabled and then the number of threads is specified. This ensures that the parallel region will have 4 threads.

```
#include <omp.h>
void main()
{
    omp_set_dynamic(0);
    omp_set_num_threads(4);
    #pragma omp parallel
    {
        int id=omp_get_thread_num();
        do_lots_of_stuff(id);
    }
}
```

# Environment Variables

- Default number of threads
  - **OMP\_NUM\_THREADS int\_literal**
- Control of dynamic mode
  - **OMP\_DYNAMIC TRUE || FALSE**
- Control of nested parallelism
  - **OMP\_NESTED TRUE || FALSE**
- Control of loop scheduling if the programmer has used `omp for schedule(RUNTIME)`
  - **OMP\_SCHEDULE "schedule[, chunk\_size]"**
- Control of threads binding
  - **OMP\_PROC\_BIND TRUE || FALSE**

---

## «Test Cases»

# Case 1: Loop & Parallel Region

- Parallelize the following sequential code with
  - parallel regions
  - worksharing

```
#define N 1024
for(int i=0; i<N; i++) { a[i] = a[i] + b[i];}
```

- OpenMP parallel region

```
#pragma omp parallel
{
    int id = omp_get_thread_num();
    int Nthrds = omp_get_num_threads();
    int istart = id * N / Nthrds;
    int iend = (id+1) * N / Nthrds;
    if (id == omp_get_num_threads()-1) iend = N;
    for(int i=istart; i<iend; i++) {a[i] = a[i] + b[i];}
}
```

adjustment for  
the last thread



# Loop & Worksharing

- Sequential code

```
#define N 1024
for(int i=0; i<N; i++) { a[i] = a[i] + b[i];}
```

- OpenMP parallel region with worksharing

```
#pragma omp parallel
{
    #pragma omp for schedule(static) default scheduling
    for(int i=0; i<N; i++) { a[i] = a[i] + b[i];}
}
```

**or simply:**

```
#pragma omp parallel for
for(int i=0; i<N; i++) { a[i] = a[i] + b[i];}
```

# Case 2: Functional parallelism

- Parallelize the following sequential code
  - what is the total execution time if each function takes one second?

```
V = alpha();  
W = beta();  
X = gamma(V, W);  
Y = delta();  
printf("%f\n", epsilon(X,Y));
```

total time = 5s

# Functional parallelism - Solution 1

```
#pragma omp parallel num_threads(3)  no sense to use more threads
#pragma omp sections
{
    #pragma omp section
    V = alpha();

    #pragma omp section
    W = beta();

    #pragma omp section
    Y = delta();
}
X = gamma(V, W);
printf("%f\n", epsilon(X, Y));
```

total time = 3s

# Functional parallelism - Solution 2

```
#pragma omp parallel num_threads(2)
{
    #pragma omp sections
    {
        #pragma omp section
        V = alpha();

        #pragma omp section
        W = beta();
    }
    #pragma omp sections
    {
        #pragma omp section
        X = gamma(V, W);

        #pragma omp section
        Y = delta();
    }
}
printf("%f\n", epsilon(X,Y));
```

no sense to use more threads

implicit barrier

total time = 3s

but with fewer threads<sup>61</sup>

# Case 3 - Reductions

- Parallelize the following sequential code

```
long num_steps = 100000;
double step;

void main ()
{
    double x, pi, sum = 0.0;
    step = 1.0/(double) num_steps;

    for (int i=0; i< num_steps; i++){
        x = (i+0.5)*step;
        sum = sum + 4.0/(1.0+x*x);
    }
    pi = step * sum;

    printf("Pi is %lf\n", pi);
}
```

# Using the reduction clause

```
long num_steps = 100000;
double step;

void main ()
{
    double x, pi, sum = 0.0;
    step = 1.0/(double) num_steps;

    #pragma omp parallel for reduction(+:sum) private(x)
    for (long i=0; i<num_steps; i++){
        x = (i+0.5)*step;
        sum = sum + 4.0/(1.0+x*x);
    }
    pi = step * sum;

    printf("Pi is %lf\n", pi);
}
```

# References

- OpenMP Specifications & Quick Reference Card
  - [www.openmp.org](http://www.openmp.org)
- OpenMP tutorial at LLNL, Blaise Barney
  - <https://computing.llnl.gov/tutorials/openMP/>
- An Overview of OpenMP, Ruud van der Pas – Sun Microsystems
  - <http://www.openmp.org/wp-content/uploads/ntu-vanderpas.pdf>