

Γνώση του A όταν οι B, C δεν γνωρίζουν το χρώμα τους

Ορισμός Έστω M_A ένα μοντέλο Kripke που αναπαριστά την γνώση του A:
Ο παίκτης A γνωρίζει ότι αληθεύει θ , αν $M_A \models \theta$.

Take any cluster S of possible worlds that A can not distinguish between.

Partition this cluster into sub-clusters, say T_1, T_2 , etc., that are indistinguishable for B.

That is, B can not distinguish between members of T_1 ; B can not distinguish between members of T_2 , and so on.

1) By [B sees A], **A knows** that either B knows $AisWh$ or B knows $\neg AisWh$.

This means that in T_1 it cannot happen that $AisWh$ is true at some worlds and false at others; similarly for T_2 , and so on.

Briefly, for each k , either $AisWh$ is true at all the possible worlds of T_k , or else $AisWh$ is false at all the possible worlds of T_k .

2) By [B hears C], **A knows** that B knows that $AisWh \vee BisWh$.

Then $AisWh \vee BisWh$ must be true throughout each T_k .

Combining with (1) it follows that, for each k :

either $AisWh$ is true at every possible world of T_k ,
or $BisWh$ is true at every possible world of T_k .

Άρα, $AisWh \vee K_B BisWh$ is true at every possible world of T_k .

Since the T_k are a partition of S , $AisWh \vee K_B BisWh$ must be true at every world of S .

Since $AisWh \vee K_B BisWh$ is true at all the worlds that A thinks are possible,

$M_A \models AisWh \vee K_B BisWh$.

3) By [B doesn't know], **A knows** that B doesn't know $BisWh$.

Then $M_A \models \neg K_B BisWh$,

and $M_A \models AisWh$.

ΕΡΩΤΗΜΑ Τι συμβαίνει όταν ο A είναι μαύρος;

Tableau proof for: $1, 2 \models (A_{isWh} \vee K_B B_{isWh})$

1 (s1) A knows that (B knows A_{isWh} or B knows $\neg A_{isWh}$)

2 (s1) A knows that B knows that ($A_{isWh} \vee B_{isWh}$)

4 (s1) $\neg (A_{isWh} \vee K_B B_{isWh})$

5 (s1) $\neg A_{isWh}$ from 4

6 (s1) $\neg K_B B_{isWh}$ from 4

7 (s1 B s3) $\neg B_{isWh}$ from 6

8 (s1) B knows that $A_{isWh} \vee B_{isWh}$ from 2

9 (s1 B s3) $A_{isWh} \vee B_{isWh}$ from 8

10 (s1 B s3) A_{isWh} from 9 , 7

11 (s1) B knows A_{isWh} **or** B knows $\neg A_{isWh}$ from 1

BRANCH 1

11; **1** (s1) B knows $\neg A_{isWh}$

12 (s1 B s3) $\neg A_{isWh}$ from 11; 1

branch closes with 12, 10

BRANCH 2

11; **2** (s1) B knows A_{isWh}

13 (s1) A_{isWh} from 11; 2

branch closes with 13, 5

Tableau proof for: $1, 2, 3 \models A \text{ knows } A \text{ isWh}$

- 1 (s1) A knows that (B knows AisWh or B knows \neg AisWh)
- 2 (s1) A knows that B knows that AisWh \vee BisWh
- 3 (s1) A knows that B doesn't know BisWh
- 4 (s1) \neg A knows AisWh

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| 5 | (s1 A s2) | $\neg \text{AisWh}$ | from 4 |
| 6 | (s1 A s2) | B doesn't know BisWh | from 3 |
| 7 | (s1 A s2) | B knows that $\text{AisWh} \vee \text{BisWh}$ | from 2 |
| 8 | (s1 A s2 B s3) | $\neg \text{BisWh}$ | from 6 |
| 9 | (s1 A s2 B s3) | $\text{AisWh} \vee \text{BisWh}$ | from 7 |
| 10 | (s1 A s2 B s3) | AisWh | from 9, 8 |
| 11 | (s1 A s2) | B knows AisWh or B knows $\neg \text{AisWh}$ | from 1 |

BRANCH 1

- 11;1 (s1 A s2) B knows \neg AisWh
- 12 (s1 A s2 B s3) \neg AisWh from 11;1

closes with 10

BRANCH 2

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|------|-----------|---------------|-----------|
| 11;2 | (s1 A s2) | B knows AisWh | |
| 13 | (s1 A s2) | AisWh | from 11;2 |

closes with 5